

Week 2. No. 1

recap: Fix a model \mathcal{M} , (\leftarrow (ctble) transitive model of)
ZFC

a poset $\mathbb{P} \in \mathcal{M}$,

a \mathcal{M} -generic $\mathcal{G} \subseteq \mathbb{P}$

Last time, we defined $\mathcal{M}[\mathcal{G}] := \{x^{\mathcal{G}} \mid x \in \mathcal{M}\}$

$x^{\mathcal{G}} := \{y^{\mathcal{G}} \mid (p, y) \in x \text{ for some } p \in \mathcal{G}\}$

showed $\mathcal{M}[\mathcal{G}]$ is transitive.

$\mathcal{M}[\mathcal{G}] \cap \text{OR} = \mathcal{M} \cap \text{OR}$

$\mathcal{M} \cup \{\mathcal{G}\} \subseteq \mathcal{M}[\mathcal{G}]$.

$\mathcal{M} \subsetneq \mathcal{M}[\mathcal{G}]$

if \mathbb{P} is atomless

$\mathcal{M}[\mathcal{G}] \models$ extensionality, foundation,
infinity, pairing, union.

To show $\mathcal{M}[\mathcal{G}] \models \text{ZFC}$,

we are missing Power set axiom,

Separation,

Replacement, and

AC (the Axiom of Choice)

All of these statements ~~are~~ are shown to be true in $M[g]$ via the forcing language

def M, \mathbb{P} fixed. Let φ be a formula in the language of set theory, ~~and~~ let $T_1, \dots, T_k \in M$ and let $p \in \mathbb{P}$.

We say p forces $\varphi(T_1 \dots T_k)$ (in \mathbb{P} over M),

written $p \Vdash_{\mathbb{P}}^M \varphi(T_1 \dots T_k)$ (or $p \Vdash \varphi(T_1 \dots T_k)$)
 (or $p \Vdash_{\mathbb{P}}^M \varphi(T_1 \dots T_k)$)

iff for every $g \in \mathbb{P}$, ~~and~~ an M -generic filter, s.t. $p \in g$, $M[g] \models \varphi(T_1^g, \dots, T_k^g)$

- The relation $p \Vdash \varphi(T_1 \dots T_k)$ is called forcing relation
- It is non-trivial that the relation $p \Vdash \varphi(T_1 \dots T_k)$ is definable inside of M

Note Fix φ .

$$\left\{ (p, T_1, \dots, T_k) \mid p \in \mathbb{P}, T_1, \dots, T_k \in M, \right. \\ \left. p \Vdash \varphi(T_1, \dots, T_k) \right\} \subseteq M^{k+1}$$

We will show that this relation is definable over M .

Lemma Fix a formula φ .

The relation

$$\left\{ (p, T_1, \dots, T_k) \mid p \in \mathbb{P}, T_1, \dots, T_k \in M, p \Vdash \frac{\mathbb{P}}{M} \varphi(T_1, \dots, T_k) \right\}$$

is definable over M

def Let $p \in \mathbb{P}$, $D \subseteq \mathbb{P}$.

We say D is dense below p if

$$\forall q \leq p \exists r \in D (r \leq q)$$

def

① Let $\tau_1, \tau_2 \in M$, let $p \in M$

$p \Vdash \tau_1 = \tau_2$ if $\forall (s_1, \pi_1) \in \tau_1$

$\left\{ q \leq p \mid q \leq s_1 \rightarrow \exists (s_2, \pi_2) \in \tau_2 \left(q \leq s_2 \wedge q \Vdash \tau_1 = \tau_2 \right) \right\}$

is dense below p , and

$\forall (s_2, \pi_2) \in \tau_2$

$\left\{ q \leq p \mid q \leq s_2 \rightarrow \exists (s_1, \pi_1) \in \tau_1 \left(q \leq s_1 \wedge q \Vdash \tau_1 = \tau_2 \right) \right\}$

is ~~dense~~ dense below p .

② Let $\tau_1, \tau_2 \in M$, let $p \in M$

$p \Vdash \tau_1 \in \tau_2$ if $\left\{ q \leq p : \exists (s, \pi) \in \tau_2 \left(q \leq s \wedge q \Vdash \tau_1 = \pi \right) \right\}$

is dense below p .

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③ Let $T_1, \dots, T_k \in M$, let $\sigma_1, \dots, \sigma_l \in M$,
 φ, ψ : formulae, $p \in \mathbb{P}$

~~$p \Vdash \varphi(T_1, \dots, T_k) \wedge \psi(\sigma_1, \dots, \sigma_l)$~~

is $p \Vdash \varphi(T_1, \dots, T_k)$ and $p \Vdash \psi(\sigma_1, \dots, \sigma_l)$

④ Let $T_1, \dots, T_k \in M$, φ a formula, $p \in \mathbb{P}$

$p \Vdash \neg \varphi(T_1, \dots, T_k)$ if for all $q \leq p$

$q \not\Vdash \varphi(T_1, \dots, T_k)$

⑤ Let $T_1, \dots, T_k \in M$, φ a formula, $p \in \mathbb{P}$

$p \Vdash \exists x \varphi(T_1, \dots, T_k)$ if

$\{q \leq p \mid \exists \sigma \in M \ q \Vdash \varphi(\sigma, T_1, \dots, T_k)\}$ is dense below p

This finishes the def.

So for each φ ,

$$\{ (p, T_1, \dots, T_k) : p \in P, T_1, \dots, T_k \in M, p \Vdash \varphi(T_1, \dots, T_k) \}$$

is definable over M

Easy lemma

For a fml φ , $T_1, \dots, T_k \in M$,

the following are equivalent:

(i) $p \Vdash \varphi(T_1, \dots, T_k)$,

(ii) $\forall q \leq p (q \Vdash \varphi(T_1, \dots, T_k))$,

(iii) $\{ q \leq p \mid q \Vdash \varphi(T_1, \dots, T_k) \}$ is dense below p .

(just looking at the definition)

Forcing theorem

Fix M , a ctble transitive model of ZFC,

$\mathbb{P} \in M$ a poset, $g \subseteq \mathbb{P}$ M -generic

Fix $\varphi, \tau_1, \dots, \tau_k \in M$.

① if $p \Vdash \varphi(\tau_1, \dots, \tau_k)$ and $p \in g$,

then $M[g] \models \varphi(\tau_1^g, \dots, \tau_k^g)$

② if $M[g] \models \varphi(\tau_1^g, \dots, \tau_k^g)$, then there is a

$p \in g$ with $p \Vdash \varphi(\tau_1, \dots, \tau_k)$

(Proof) By induction, we simultaneous^a prove
(1) + (2).

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① for φ being an identity statement:

let $p \in \mathbb{Q}$, $p \Vdash T_1 = T_2$.

We show $T_1^q \subseteq T_2^q$ (by symmetry, also get $T_2^q \subseteq T_1^q$)

Let $x \in T_1^q$, say $x = \pi_1^q$, where $(s_1, \pi_1) \in T_1, s_1 \in \mathbb{Q}$.
(Need to see: $x \in T_2^q$)

Pick $r \in \mathbb{Q}$ with $r \leq p, s_1$.

We have $r \Vdash T_1 = T_2$, since $p \Vdash T_1 = T_2$.

Hence by def, we have

$$D = \left\{ r \leq r \mid r \leq s_1 \rightarrow \exists (s_2, \pi_2) \in T_2 (r \leq s_2 \wedge r \Vdash \pi_1 = \pi_2) \right\}$$

is dense below $r \in \mathbb{Q}$.

Look at $D' = D \cup \{s \in \mathbb{R} \mid \text{there is no } t \in \mathbb{R} \text{ } t \leq s, r\}$

Then D' is dense in \mathbb{R} , $D' \in M$.

Since g is M -generic, $D' \cap g \neq \emptyset$.

Let $s \in D' \cap g$. Since $r \in \mathbb{Q}$, $s \in D$.

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Since $S \leq S_1$, there is $(S_2, \pi_2) \in T_2$ s.t.

$$S \leq S_2 \text{ and } S \Vdash \pi_1 = \pi_2$$

Then by induction, $\pi_1^g = \pi_2^g$.

Therefore $\chi = \pi_1^g = \pi_2^g \in T_2^g$ (as $S_2 \in g$).

We proved $\pi_1^g \in \pi_2^g$.

By symmetry, $\pi_1^g = \pi_2^g //$

② for φ being an identity:

Let $T_1^g = T_2^g$. Consider:

$$\psi_1(r) \equiv \exists (s_1, \pi_1) \in T_1 \left(r \leq S_1 \wedge \forall (s_2, \pi_2) \in T_2 \forall g \right. \\ \left. (g \leq S_2 \wedge g \Vdash \pi_1 = \pi_2 \rightarrow g \perp r) \right)$$

$\neg \exists t (t \leq g, r)$
 g, r are incompatible

Exercise

① if $r \in g$, then $\neg \psi_1(r)$ ② if $r \in g$, then $\neg \psi_2(r)$

$$\psi_2(r) \equiv \exists (s_2, \pi_2) \in T_2 \left(r \leq S_2 \wedge \forall (s_1, \pi_1) \in T_1 \forall g \left(r \leq S_1 \wedge g \Vdash \pi_1 = \pi_2 \rightarrow g \perp r \right) \right)$$

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Proof of exercise We only check ①.

Assume we had $\varphi_1(r)$ for some $r \in \mathbb{Q}$.

Let $(s_1, \pi_1) \in \mathbb{R}$ witness this fact.

Let $(s_2, \pi_2) \in \mathbb{T}_2$ be s.t. $s_2 \in \mathbb{Q}$ and $\pi_1^{\mathbb{Q}} = \pi_2^{\mathbb{Q}}$.

~~Let $q \in \mathbb{Q}$ be s.t. $q \leq$~~ By induction, there is $q_0 \in \mathbb{Q}$
s.t. $q_0 \vdash \pi_1 = \pi_2$. Let $q \in \mathbb{Q}$ be s.t. $q \leq q_0, s_2$.

Then we must have $q \perp r$.

But $q, r \in \mathbb{Q}$. This is a contradiction !!

$$\text{Let } D = \left\{ r \in \mathbb{R} \mid \varphi_1(r) \vee \varphi_2(r) \vee r \vdash \pi_1 = \pi_2 \right\}$$

ex

D is dense

proof of ex

Given any $p \in \mathbb{P}$. Assume $p \nmid T_1 = T_2$

Then there is some $(\pi_1, s_1) \in T_1$ s.t.

$$D = \left\{ q \leq p \mid q \leq s_1 \rightarrow \exists (\pi_2, s_2) \in T_2 (q \leq s_2 \wedge q \nmid \pi_1 = \pi_2) \right\}$$

is not dense below p . (By symmetry, we may assume)

Then there is $q \leq p$ s.t. $\forall r \leq q (r \notin D)$

Then $\psi_1(q)$ holds //
(and $q \leq p$)

Therefore, $p \in q \cap D$.

By ex, we cannot have $\psi_1(p)$, neither $\psi_2(p)$

Hence $p \nmid T_1 = T_2$ //

① for $T_1 \in T_2$: assume $q \ni p \vdash T_1 \in T_2$.

Hence there is $q \in \mathcal{Q}$ s.t.

$$q \leq p \wedge \exists (s, \pi) \in T_2 (q \leq s \wedge q \vdash T_1 = \pi)$$

Say $(s, \pi) \in T_2$ witnesses that.

Then by induction, $T_1^q = \pi^q$.

Since $q \leq s$ and $q \in \mathcal{Q}$, $s \in \mathcal{Q}$.

Hence $T_1^q = \pi^q \in \pi^q$ ✓

② for $T_1 \in T_2$: Say $T_1^q \in T_2^q$

Then there is $(s, \pi) \in T_2$ s.t.

$$s \in \mathcal{Q} \text{ and } T_1^q = \pi^q.$$

Then we have $q \ni r \vdash T_1 = \pi$.

Let $p \in \mathcal{Q}$, $p \leq s, r$

Then $p \vdash T_1 \in T_2$ by def. ✓

①, ② for $\varphi \wedge \psi$: Easy.

① for negation:

$$p \in \mathcal{Q}, p \vdash \neg \varphi(\tau_1, \dots, \tau_k)$$

~~Suppose~~ Suppose $M[\mathcal{Q}] \models \varphi(\tau_1^{\mathcal{Q}}, \dots, \tau_k^{\mathcal{Q}})$,

Then by induction, $\exists q \in \mathcal{Q}$ s.t.

$$q \vdash \varphi(\tau_1, \dots, \tau_k).$$

Let $r \leq q, p$ with $r \in \mathcal{Q}$.

~~Then $r \vdash \varphi(\tau_1, \dots, \tau_k)$ and $r \leq p$.~~

Then $r \vdash \varphi(\tau_1, \dots, \tau_k)$ and $r \leq p$.

This is a contradiction \llcorner \llcorner

② for negation:

$$\text{Suppose } M[\mathcal{Q}] \models \neg \varphi(\tau_1^{\mathcal{Q}}, \dots, \tau_k^{\mathcal{Q}})$$

$$D = \{ p \in \mathcal{Q} \mid p \Vdash \varphi(t_1, \dots, t_k) \text{ or } p \Vdash \neg \varphi(t_1, \dots, t_k) \}$$

is dense. (by def)

Let $p \in D \cap q$

Then we must have $p \Vdash \neg \varphi(t_1, \dots, t_k)$ /

① for existential statements:

$$q \Vdash p \Vdash \exists x \varphi(x, t_1, \dots, t_k)$$

By def, we have $q \in \mathcal{Q}$ s.t.

$$q \leq p \wedge \exists b \in M \quad q \Vdash \varphi(b, t_1, \dots, t_k)$$

Then by induction

$$M[q] \Vdash \varphi(b^q, t_1^q, \dots, t_k^q) \quad /$$

$$(\rightarrow M[q] \Vdash \exists y \varphi(y, t_1^q, \dots, t_k^q))$$

② for existential^{ia} statements:


Assume $M \models \exists x \varphi(x, t_1, \dots, t_k)$

$\sigma \in M$ s.t.

$$M \models \varphi(\sigma, t_1, \dots, t_k)$$

By induction, we have $p \in \mathcal{Q}$ s.t.

$$p \Vdash \varphi(\sigma, t_1, \dots, t_k). \quad \text{ok} //$$
$$(\rightarrow p \Vdash \exists x \varphi(x, t_1, \dots, t_k))$$

This finishes the proof of thm 

We keep fixing a cttle transitive model M of ZFC, No. 16
 a poset $\mathbb{P} \in M$,
 (sometimes) a M -generic filter $\mathcal{g} \subseteq \mathbb{P}$.

def

$$\begin{array}{c}
 \mathbb{P} \Vdash \mathbb{P} \\
 \uparrow \\
 \mathbb{P}
 \end{array}
 \Vdash_M \underbrace{\varphi(\tau_1, \dots, \tau_k)}_{\uparrow M} \text{ iff for all } h \subseteq \mathbb{P} \text{ } M\text{-generic} \\
 \text{with } p \in h, \\
 M[h] \models \varphi(\tau_1^h, \dots, \tau_k^h)$$

We also defined

$$p \Vdash \varphi(\tau_1, \dots, \tau_k) \quad (\text{syntactical version})$$

We showed:

Forcing theorem, part 1

① if $p \in \mathcal{g}$, $p \Vdash \varphi(\tau_1, \dots, \tau_k)$,

then $M[\mathcal{g}] \models \varphi(\tau_1^{\mathcal{g}}, \dots, \tau_k^{\mathcal{g}})$

② if $M[\mathcal{g}] \models \varphi(\tau_1^{\mathcal{g}}, \dots, \tau_k^{\mathcal{g}})$, then

there is some $p \in \mathcal{g}$ s.t. $p \Vdash \varphi(\tau_1, \dots, \tau_k)$

2nd part of the Forcing theoremLet φ fml, $\tau_1, \dots, \tau_k \in M$ ① for all $p \in \mathbb{P}$

$$p \Vdash_{\mathbb{P}}^M \varphi(\tau_1, \dots, \tau_k) \text{ iff } p \Vdash \varphi(\tau_1, \dots, \tau_k)$$

② If $M[g] \models \varphi(\tau_1^g, \dots, \tau_k^g)$, then there is $p \in g$

$$\text{s.t. } p \Vdash_{\mathbb{P}}^M \varphi(\tau_1, \dots, \tau_k)$$

(proof)

① (\Rightarrow)Suppose $p \Vdash \varphi(\tau_1, \dots, \tau_k)$.

Suppose that

$$p \not\Vdash \varphi(\tau_1, \dots, \tau_k)$$

Let $q \leq p$ be s.t.for all $r \leq q$

$$r \not\Vdash \varphi(\tau_1, \dots, \tau_k)$$

Then ~~we~~ by def,we have $q \Vdash \neg \varphi(\tau_1, \dots, \tau_k)$ Let $h \subseteq \mathbb{P}$ be M -generic with $q \in h$. Then $M[h] \models \neg \varphi(\tau_1^h, \dots, \tau_k^h)$ But since $p \in h$, we have $M[h] \models \varphi(\tau_1^h, \dots, \tau_k^h)$.

This is a contradiction !!

- | ①

(\Leftarrow) Given h M -generic with $p \in h$.Then the Forcing theorem
(by) part 1, ①,

we have

$$M[g] \models \varphi(\tau_1^h, \dots, \tau_k^h)$$

 \rightarrow we have

$$p \Vdash \varphi(\tau_1, \dots, \tau_k) //$$

② : clear //

→ thm

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Theorem

$M[g] \models ZFC$

(proof)

We checked

- Infinity
- Pairing
- Foundation
- Extensionality
- Union

Separation

We need to see :

in $M[g]$: if $x, z_1, \dots, z_k \in M[g]$ and

φ is a fml, then

$\{y \in x \mid \varphi(y, z_1, \dots, z_k)\}$ exists

Say $\tau^g = x, \tau_1^g = z_1, \dots, \tau_k^g = z_k$.

Let

$$\pi = \left\{ (p, p) : \exists q \geq p \left((q, p) \in \tau, p \Vdash_M \varphi(\overset{p}{\cancel{p}}, \tau_1, \dots, \tau_k) \right) \right\}$$

$\in M$ by the Forcing theorem $\left(\begin{array}{l} p \Vdash \varphi(p, \tau_1, \dots, \tau_k) \\ \text{is definable over } M \end{array} \right)$



Want

$$\pi^g = \left\{ y \in X \mid \varphi(y, z_1, \dots, z_k) \right\}$$

" \subseteq " Let $(p, p) \in \pi$ with $p \in g$.

Then by def of π ,

$p^g \in \tau^g = X$, since for some $q \geq p$, $(q, p) \in \tau$,

and $M[g] \models \varphi(p^g, \tau_1^g, \dots, \tau_k^g)$, since

$p \Vdash \varphi(p, \tau_1, \dots, \tau_k)$.



$M[g] \models$

OK

" \supseteq " Given $y \in X$ w/ $\varphi(y, z_1, \dots, z_k)$. ~~Say~~

~~Let $q \in g$ be s.t.~~

Let $(q, p) \in \tau$ be s.t. $q \in g$ and $p^g = y$.

By the Forcing theorem, there is $p' \in g$ s.t.

$p' \Vdash \varphi(p, \tau_1, \dots, \tau_k)$.

Let $p \in g$ s.t. $p \leq q, p'$. (g : filter)

Then by def of π , $(p, p) \in \pi$.

Hence $p^g \in \pi^g$, since $p \in g$ // ok

This finishes the proof of Separation.

Replacement

Assume $M[g] \models \forall x \in T^g \exists y \varphi(x, y, T_1^g, \dots, T_k^g)$

We need to see $\exists z \in M[g]$ s.t.

$$M[g] \models \forall x \in T^g \exists y \in z \varphi(x, y, T_1^g, \dots, T_k^g)$$

$$\text{Let } \pi = \left\{ (p, \sigma) : \exists q \geq p \exists \bar{\sigma} \left((q, \bar{\sigma}) \in T \wedge p \Vdash \varphi(\bar{\sigma}, \sigma, T_1, \dots, T_k) \right) \right\}$$

π seems to work, but π is not a set
→ we need to restrict σ

Redefine, fix suff large $\alpha \in \text{OR} \cap M$, $\alpha < \text{ON} \cap M$

$$\pi = \left\{ (p, \sigma) : \exists q \geq p \exists \bar{\sigma} \left((q, \bar{\sigma}) \in T \wedge p \Vdash \varphi(\bar{\sigma}, \sigma, T_1, \dots, T_k) \wedge \sigma \in V_\alpha^M \right) \right\}$$

Then π is a set in M !!

How to find such an α ?

Use Replacement in M !

Consider the function.

$F: (p, \delta) \mapsto$ the least $\alpha \in ON \cap M$ s.t.

$$\exists q \geq p \exists \bar{\delta} \text{ s.t. } (q, \bar{\delta}) \in T \wedge$$

$$p \Vdash \varphi(\bar{\delta}, \delta, T_1, \dots, T_k) \wedge$$

$$\delta \in V_\alpha^M$$

Then by replacement in M , $\text{ran}(F)$ is a set in M .

Let $\alpha = \sup \text{ran}(F) \in ON \cap M$

Then α works //

And π^g is a desired set // (same argument as for separation.)

Power set Axiom

Exercise! (same)

$$\begin{aligned} V_\alpha &= \{ x : rk_\epsilon(x) < \alpha \} \\ &= \bigcup \{ P(V_\beta) \mid \beta < \alpha \} \end{aligned}$$

$$V_0 = \emptyset$$

$$V_{\alpha+1} = P(V_\alpha)$$

$$V_\alpha = \bigcup_{\beta < \alpha} V_\beta, \text{ if } \alpha \text{ is limit}$$

We want to see that if $\chi \in M[g]$,

then

$$M[g] \models \exists \alpha \exists \text{ injection } f: \chi \longleftrightarrow \alpha$$

Fix $\chi = \tau^g \in M[g]$.

In M , there is some α and some

bijection $\bar{f}: \tau \rightarrow \alpha$

● We may inside M assign to every pair σ, π some $\langle \sigma, \pi \rangle$ s.t. $\langle \sigma, \pi \rangle^g = (\sigma^g, \pi^g)$

Let

$$\tau = \left\{ (p, \langle p, \xi \rangle) \mid \begin{array}{l} \text{[Scribbled out]} \\ (p, p) \in \tau \wedge \bar{f}((p, p)) = \xi \end{array} \right\}$$

τ^g is 1-to-many relation.

So we can make τ^g to an injective function by separation in $M[g]$

(just minimize the range)

We showed $M[g] \models \text{ZFC}$ //



We developed the theory of forcing \dagger

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We showed $M[G] \models ZFC$

Next: Vary Cohen forcing to produce a

$$M[G] \models 2^{\aleph_0} = \aleph_2$$

Before doing this, we'll produce a

$$M[G] \models \underbrace{2^{\aleph_0} = \aleph_1}_{CH}$$