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Motivation: Want to compare hod pairs
(P, Σ), (Q, Λ)

problem: copying Σ, Λ.

want to reach $i: P \xrightarrow{I} P^*$ via Σ
 $j: Q \xrightarrow{u} Q^*$ via Λ

s.t. $P^* \sqsubseteq Q^*$ or $Q^* \sqsubseteq P^*$ and

the tails $\sum_{I, P^*}, \wedge_{u, Q^*}$ are
copyable.

convention, our strategies act on stacks \vec{I}
of normal trees: \vec{I}_0 on base model,
 $\vec{I}_{\alpha+1}$ on last model of \vec{I}_α , \vec{I}_λ on
dir. lin of \vec{I}_α , $\alpha < \lambda$.

normal: lengths of extenders increasing +

if E is used before F on the
same track, then $crit(F) \geq \lambda_E = i_E(crit(E))$

$$(w \nu(E) \leq crit(F))$$

$$\sum_{\vec{I}} (\vec{u}) = \sum (\vec{I} \wedge \vec{u})$$

prob: iterating away strategy disagreements...
how to do it?

$$\left. \begin{matrix} \vec{I} \\ P_\alpha \end{matrix} \right\} \left. \begin{matrix} \vec{u} \\ Q_\alpha \end{matrix} \right\} \text{ say } \sum_{\vec{I}, R} \neq \sum_{\vec{u}, R}$$

idea: "use brack embeddng as embed"

problem: more generators!!

if the 2 are no ext. overlapping woods in
the had mice, then the 2 ways are
by woods, Sargoyan.

but not in general.

potential solution: compare (P, Σ) and
 (Q, Λ) with a case beyond control.

(*) (P, Σ) : for any suff. good
background universe caply (P, Σ) ,
 (P, Σ) iterates into its $[E \vec{I}]$
or its had mouse constr., no strategy

disagreement, no extended disagreement.

here: "copy" means the following.

in $B =$ background universe

we have: $P \in HC^B$, Σ has u.b.
code in B .

this proof needs condensation properties for Σ . (to construct Σ with these properties we need some form of UBH in the background universe when it is constructed at (P, Σ) .)

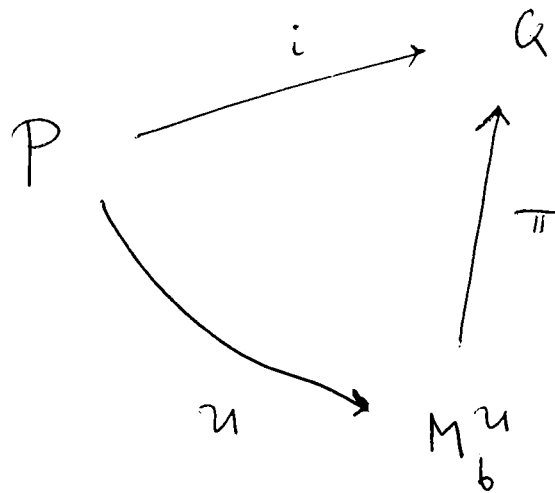
Sample condensation properties:

(1) hull condensation.

if \mathcal{I} is by Σ and \mathcal{U} is a hull of \mathcal{I} , then \mathcal{U} is by Σ .

Sargis showed UBH yields such (P, Σ) .

(2) branch condensation :



i by Σ ,
 $\pi \circ i_b^u = i$
 u by Σ ,
 then
 $\Sigma(u) = b$.

e.g. M is a free module, $\rho_w(M) = w$,
 M sub, M ideal.

then M ~~sub~~ has a unique Σ ,
 $\Sigma(u) = \text{unique } b \text{ s.t. } M_b^u \text{ is ideal.}$

outline of $(*) (P, \Sigma)$, when Σ has
 branch condensation.

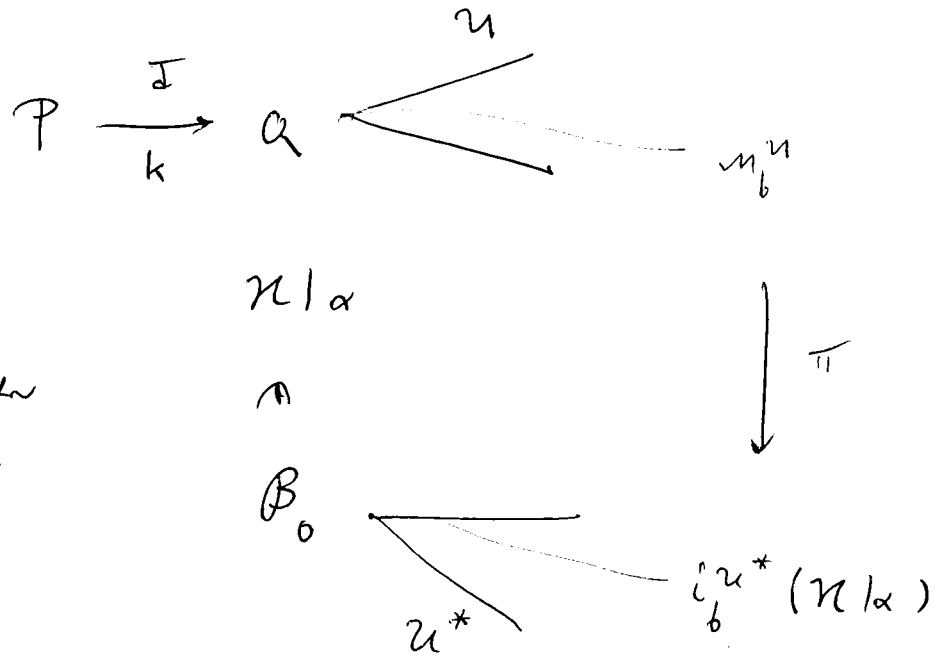
\mathcal{K} some level of the $L[\vec{E}]$ (or local
 monoid).

(P, Σ) , $P \in \mathcal{K}$, Σ u. b.

$$P \longrightarrow Q$$

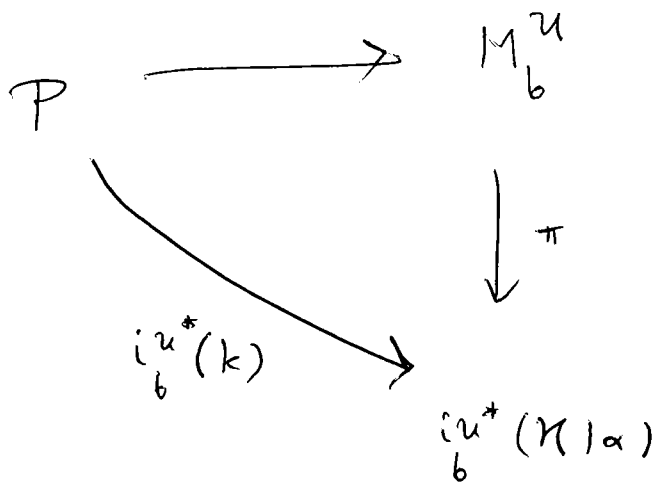
say $Q|\alpha = \mathcal{K}|\alpha$, $\sum_{\vec{d}, a|\alpha} (u) \neq \Omega(u)$,

where Ω is the strategy for π .



B_0 = background for the construction of π

$i_b^{u^*}(I)$ has been noted $i_b^{u^*}(K|\alpha)$, and $i_b^{u^*}(I)$ is by Σ (because Σ is universally bair)



hence b is acc. to Σ , so $b = \sum_{\mathcal{I}} (u)$. \dashv

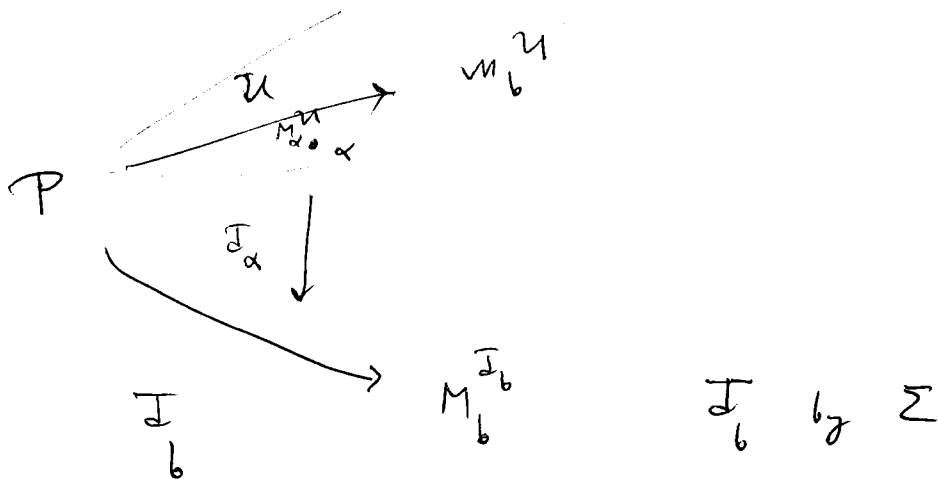
problem : brach condensation probably fails once we put to exclude overlapping worldlines.

($\Sigma(\mathcal{I}) \neq$ the map c s.t. $M_c^{\mathcal{I}}$ is isom.

on γ . $\Sigma(\mathcal{I}) =$ map b s.t.

$\Phi(\mathcal{I} \hat{=} b)$ is isom.

phalax condensation



possibly, background condensation yielded (P, Σ) with phalax condensation.

can show $(*) (P, \Sigma)$ if

Σ has phalanx condensation

~~the~~ above condensation property :

normalizing cell : $\mu_\alpha (\bar{I}, u)$, a

2-stack of normal trees,

can construct a normal tree $W(\bar{I}, u)$,

minimal in some sense, s.t. every

$$M_\alpha^u \xrightarrow{\pi_\alpha} M_{\gamma(\alpha)}^{W(\bar{I}, u)}$$

—
 $\ln W(\bar{I}, u \hat{=} b) = W_b$

in our branch condensation app,

W_b is a pseudo hull of $i_b^*(\bar{I})$.

our Σ will have strong hull condensation

which means W_b is by Σ .

def. Σ 2-normalized well iff
 whenever (I, u) is by Σ , then
 $W(I, u)$ is by Σ

so if $c = \sum_I(u) = \sum(I \wedge u)$,
 then W_c is by Σ .

can show that this implies ~~by~~

def. $W(I, u)$
 normalizing well
 by full condensation
 background construction
 lifting iterability

fine structure : do it for pure ext. premice
 λ -indexing .

convention. every premouse M has a
associated $k(M)$ (its degree of
soundness)

$$\text{ult}(M; E) = \text{ult}_{k(M)}(M; E)$$

gives rise to Σ_{k+1} el. embedding .

$$\rho(M) = \rho_{k(M)+1}(M) .$$

let \mathcal{I} be a normal tree on M ,
last model Q .

let F be on the Q -sequence .

define $W(\mathcal{I}, F)$

let $\alpha = \alpha^{\mathcal{I}, F}$ be least s.t .

F is on the $M_\alpha^{\mathcal{I}}$ -sequence .

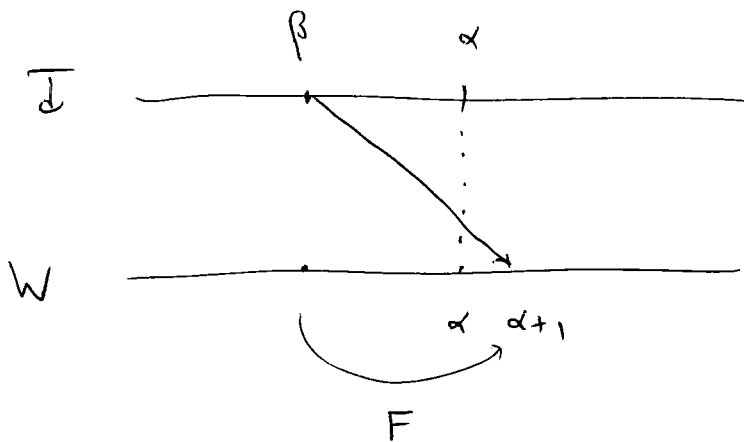
let $\beta = \beta^{\mathcal{I}, F}$ = least β s.t .

$$\mu = \text{crit}(F) < \lambda_{E_\beta^{\mathcal{I}}} .$$

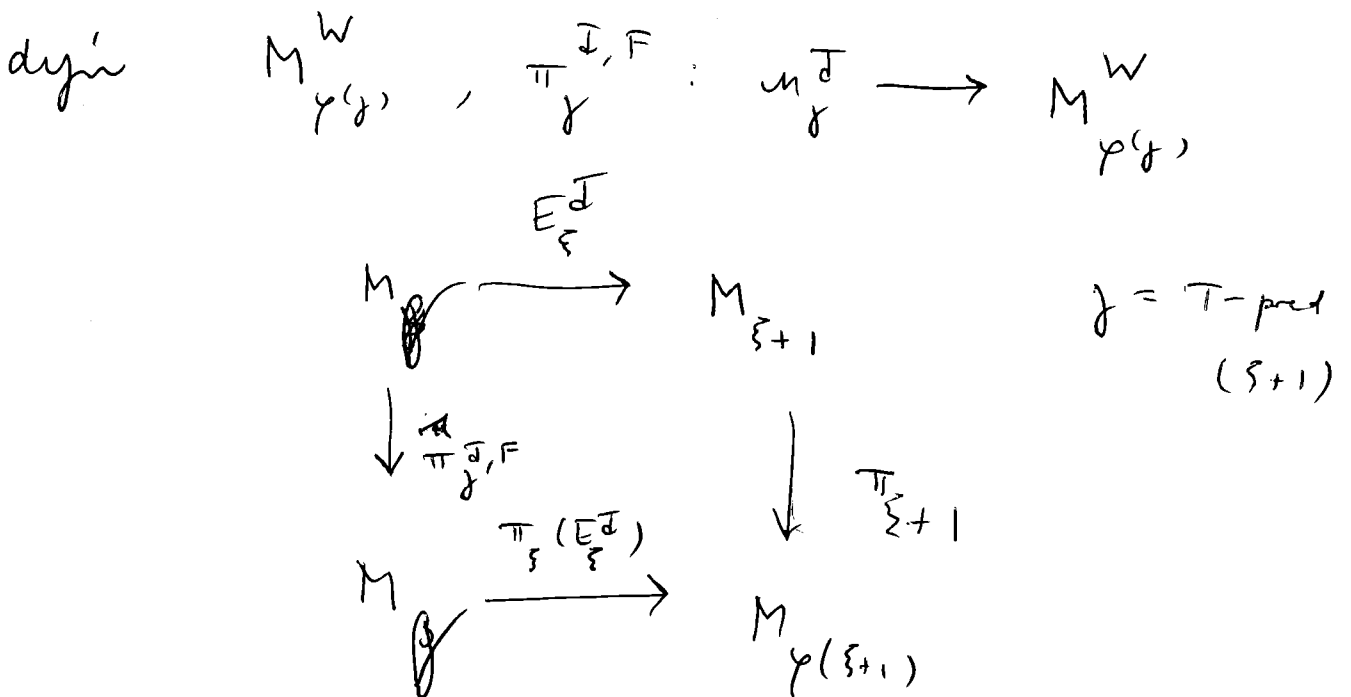
$$W(\mathbb{I}, F) \upharpoonright (\alpha^{\mathbb{I}, F} + 1) = \mathbb{I} \upharpoonright (\alpha^{\mathbb{I}, F} + 1)$$

$$M_{\alpha+1}^W = \text{ult}(P, F)$$

$P \trianglelefteq M_{\beta}^{\mathbb{I}}$, dictated by nonality.



remains part of W via by copying.



for $\text{cnt}(E_\xi^{\mathbb{J}}) < \lambda_{\mathbb{J}}^\beta = \sup \{ \lambda_{E_\alpha^{\mathbb{J}}} : \alpha < \beta \}$

or

$\text{cnt}(E_\xi^{\mathbb{J}}) \geq \mu$

o.w. $\beta = \text{T-pred}(\xi_{+1})$

$$\begin{array}{ccc}
 M_\beta & \xrightarrow{E_\xi} & M_{\xi_{+1}} \\
 \downarrow \text{id} & & \downarrow \pi_{\xi_{+1}} \\
 M_\beta & \xrightarrow{\pi_\xi(E_\xi)} & M_{\rho(\xi_{+1})}^W
 \end{array}$$

$$\varphi(\gamma) = \begin{cases} \gamma & \text{if } \gamma < \beta \\ (\alpha+1) + (\gamma - \beta) & \text{if } \gamma \geq \beta \end{cases}$$

can we

$$W(\mathbb{J}, F) = \mathbb{J}^{< \text{rk}(F)} \frown \langle F \rangle \frown \underset{F}{\text{''}} \mathbb{J}^{> \text{cnt}(F)}$$

branch / extend

\mathcal{I} normal on \mathcal{P}
 u on last node of \mathcal{I} .

$$W(\mathcal{I}, u)$$

notice F could have come from m_α^S , where S is a normal tree on \mathcal{P} .

say $\beta = \beta^{S, F}$, $\alpha = \alpha^{S, F}$.

as long as \mathcal{I} on \mathcal{P} is normal w/

$$M_\beta^{\mathcal{I}} = M_\beta^S$$

(so in fact

$$\mathcal{I} \upharpoonright \beta+1 = S \upharpoonright \beta+1$$

the if \mathcal{I}, S by same strategy)

can make then γ $W(\mathcal{I}, F)$.

$$W(\mathcal{I}, F) = S \upharpoonright \alpha+1 \hat{\langle F \rangle} \hat{i}_F \text{'' } \mathcal{I} \upharpoonright \gamma \text{'' } \text{ent}(F)$$

now define $W(\mathcal{I}, u)$ by ind. on γ

$$W_\gamma = W(\mathcal{I}, u \upharpoonright \gamma+1)$$

$$W_0 = \mathbb{I}.$$

similarly, defi, for $\nu < \omega$ &

$$\gamma_{\nu, \delta} : \text{el } W_\nu \longrightarrow \text{el } W_\delta$$

(possibly partial,
happy if

$(\nu, \delta]_\omega$ has a drop).

$$\text{then } \pi_{\tau}^{\nu, \delta} : M_{\tau}^{W_\nu} \longrightarrow M_{\tau}^{W_\delta} \\ \gamma_{\nu, \delta}(\tau)$$

for $\tau \in \text{dom } \gamma_{\nu, \delta}(\tau)$.

these have various agreement properties.

$\gamma_{\nu, \delta}, (\pi_{\tau}^{\nu, \delta} : \tau \in \text{dom})$ is a

example of a pseudo hull

then also: let $R_\delta = \text{last model of } W_\delta$

$$\text{then } \sigma_\delta : M_\delta^u \longrightarrow R_\delta$$

set $F_j = \sigma_j(E_j^u)$.

$$W_{j+1} = W(W_\nu, F_j),$$

whv $\nu = u - \text{pred}(j+1)$.

$$\beta^{W_\nu, F_j}$$

$\alpha_j = \text{lean } \alpha \text{ s.t. } F_j \text{ is on the } M_\alpha^{W_j} \text{-set.}$

so $W_{j+1} \uparrow(\alpha_{j+1}) = W_j \uparrow(\alpha_j + 1)$.

$$\gamma_{\nu, j+1}(\xi) = \begin{cases} \xi & \text{if } \xi < \beta^{W_\nu, F_j} \\ (\alpha_j + 1) + (\xi - \beta) & \text{if } \xi \geq \beta \end{cases}$$

defined this way if $\nu \mapsto j+1$ does not drop in model or degree.

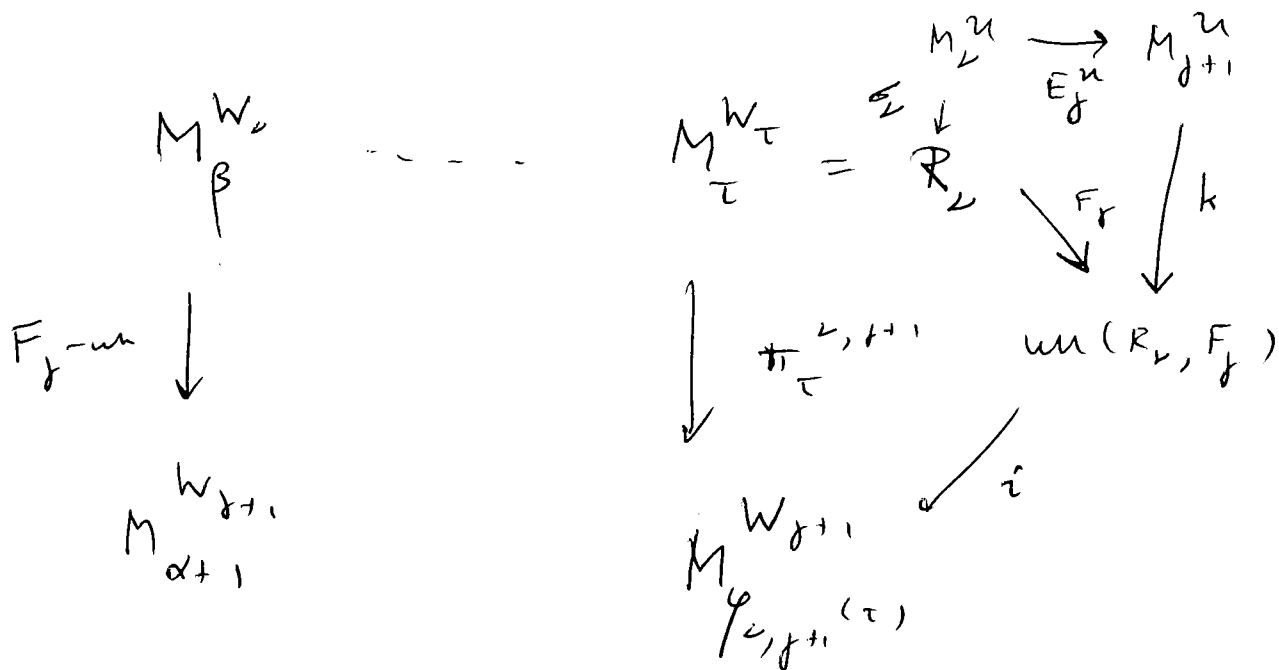
o.w. $\text{dom } \gamma_{\alpha, \delta+1} = \beta + 1$.

$$\pi_{\tau}^{\alpha, \delta+1} : M_{\tau}^{W_{\alpha}} \longrightarrow M_{\gamma_{\alpha, \delta+1}(\tau)}^{W_{\delta+1}}$$

$$\gamma_{\gamma, \delta+1} \quad \text{or } \gamma <_{\alpha} \gamma$$

$$\gamma_{\alpha, \delta+1} \circ \gamma_{\gamma, \delta}$$

$$\pi_{\tau}^{\gamma, \delta+1} = \pi_{\tau}^{\alpha, \delta+1} \circ \pi_{\tau}^{\gamma, \alpha}$$



$$\sigma_{\alpha} \uparrow \text{un } E_{\alpha}^u = \sigma_{\delta} \uparrow \text{un } E_{\delta}^u$$

$$\sigma_{\delta+1} = i \circ k$$